

# SIMULATED ANNEALING

## A METHODE TO SOLVE COMBINATORIAL PROBLEMS

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# OUTLINE

- 1 COMBINATORIAL PROBLEMS
- 2 THE ALGORITHM
- 3 PARAMETER
- 4 CONCLUSION AND SOURCES

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- 1 COMBINATORIAL PROBLEMS
- 2 The Algorithm
- 3 Parameter
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# WHAT IS A COMBINATORIAL PROBLEM?

- Which problems
- Parameters
- Classic examples
- given  $n$  objects to observe
- search for optimum
- Focus: "better" result

# WHAT IS A COMBINATORIAL PROBLEM?

- Which problems
- **Parameters**
- Classic examples
- $S$ : solutions space
- $f$ : cost function
- $f(i)$ : quality of solution

# WHAT IS A COMBINATORIAL PROBLEM?

- Which problems
- Parameters
- Classic examples
- Clique
- TSP
- Hamilton-Path

# RESOLUTION METHODES

- Numerical  
methode
- Heuristical  
methode
- "brute force"
- searching in the whole  $S$
- pro: find the global optimum
- con: huge complexity

# RESOLUTION METHODES

- Numerical methode
- Heuristical methode
- stochastic
- local search:
- iteration with limited workspace
- searching in the neighbourhood

# LOCAL SEARCH

## HOW DOES IT WORK?

### INITIALIZATION

Choose a random solution  $i$  from the workspace (TSP: choose a random permutation for the trip)

### ITERATION

1. Try to find a better solution on the neighbourhood of  $i$  (TSP: change the route a bit)
2. Record better solutions
3. Iterate back
4. Terminate if the whole neighbourhood is completely searched

# LOCAL SEARCH

## HOW DOES IT WORK?

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### ITERATION

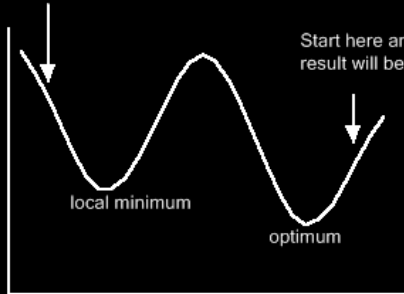
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# RESULT

## UNRELIABLE

The result is often not the optimum one, it depends a lot on the initialization.

Start here and the  
result will be local minimum



# WEAKNESS OF LOCAL SEARCH

## PROBLEM

The algorithm only finds a local minimum

## CAUSES

1. Only better results are accepted
2. The local minimum cannot be overcome

## IMPROVEMENT

How about enlarging the neighbourhood of  $i$ ?

## WHY NOT?

No way!! Since otherwise it will look like numerical methods

## OTHER POSSIBILITY

Also accept worse solutions

## BUT..

1. How far we will accept them?
2. Which rules do we use?

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# BACKGROUND

## PHYSICAL PHENOMENA AS THE BASIS

### METROPOLIS [1953]

Experiment on physical objects (liquids, gasses, solids) of their temperature, pressure, etc.

### ANNEALING PROCESS OF LIQUIDS

Release a minimal amount of energy

### SLOW ANNEALING

Stable and steady structure (minimal energy balance)

### FAST ANNEALING

Unsteady structure, unbalance energy release (non-optimal)

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# METROPOLIS

## PART I

### IMPORTANT

To overcome local minima we might also accept some worse solutions

### PROBABILITY

We use a distribution which dependent on

- Temperature  $T$  (mobility of the particles)
- Energy difference  $\Delta E = E_2 - E_1$

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# METROPOLIS

## PART II

Let  $p$  be acceptance probability

### ALGORITHM

| better solution    | worse solution                                                              |
|--------------------|-----------------------------------------------------------------------------|
| $E_2 < E_1$        | $E_2 > E_1$                                                                 |
| accept ( $p = 1$ ) | accept with $p = e^{-\frac{E_2 - E_1}{T}}$ and $p \geq \text{random}[0; 1]$ |

# METROPOLIS

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# THE GENERAL ALGORITHM

## PRINCIPLE PROCEDURE

### PSEUDO-CODE

```
INIT(i_start, T_start, HT_start);  
k:=0; i:=i_start; T_k:=T_start;  
repeat  
  for h:=1 to HT_k  
    GENERATE(j of S_i);  
    if  $f(j) < f(i)$  then  $i:=j$ ;  
    else  
      if  $e^{(f(i)-f(j))/T_k} > \text{random}[0;1]$  then  $i:=j$ ;  
  k:=k+1;  
  DETERMINE_NEIGHBOURHOOD_SIZE(HT_k);  
  SCHEDULE_ANNEALING(T_k);  
until End_Conditions reached;
```

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# PARAMETER SELECTION

- Initial temperature
- Annealing schedule
- Neighbourhood  $H_T$
- End Conditions
- Known issues

## SETTING

- Must be high enough to ensure result independent from initial solution

# PARAMETER SELECTION

- Initial temperature
- **Annealing schedule**
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## SETTING

- geometrical:  $d(T) = c \cdot T$  where  $c \in [0, 8; 0, 99]$
- slow annealing:  $d(T) = \frac{T}{1+c \cdot T}$  where  $c > 0$  and small

# PARAMETER SELECTION

- Initial temperature
- Annealing schedule
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## SETTING

- $H_T$  must be reasonable large
- large  $H_T$  at initial stages is unnecessary since it will be "neutralized" by the initial random solution
- intensive searching towards the end of algorithm is needed to reach optimum

⇒ Increase  $H_T$  by temperature reduction

# PARAMETER SELECTION

- Initial temperature
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## SETTING

Combination of:

- reaching an end-temperature or exceeding a maximal number of iterations
- at some large interval of runtime no better solution found
- acceptance rate falls below a fixed limit (e.g. under 3%)

# PARAMETER SELECTION

- Initial temperature
- Annealing schedule
- Neighbourhood  $H_T$
- End Conditions
- **Known issues**

## SETTING

- parameter setting is the significant part to the algorithm
- SA works fine only if the parameters set correctly
- too low temperature: get stuck on a local minimum
- too high temperature: the optimum will be passed several times
- too large neighbourhood: high complexity

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## SIMULATED ANNEALING

Meta - heuristical optimization methodes to solve combinatorial problems

### BASIC

On physical process in the nature: annealing material

### IDEA

Accept worse solution to gain the optimum one

### EVALUATION

- + easy to implement
- + good end result (1-10% to optimum)
- + several hundreds possible application
- difficult parameter selection

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# LITERATURES

- “Simulated Annealing” - im Rahmen des PS Virtual Lab” Martin Pfeiffer.
- Simulated annealing - Wikipedia, the free encyclopedia. Online. 01 May 2007.  
<[http://en.wikipedia.org/wiki/Simulated\\_annealing](http://en.wikipedia.org/wiki/Simulated_annealing)>
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<<http://www-i1.informatik.rwth-aachen.de/~algorithmus/algo41.php>>